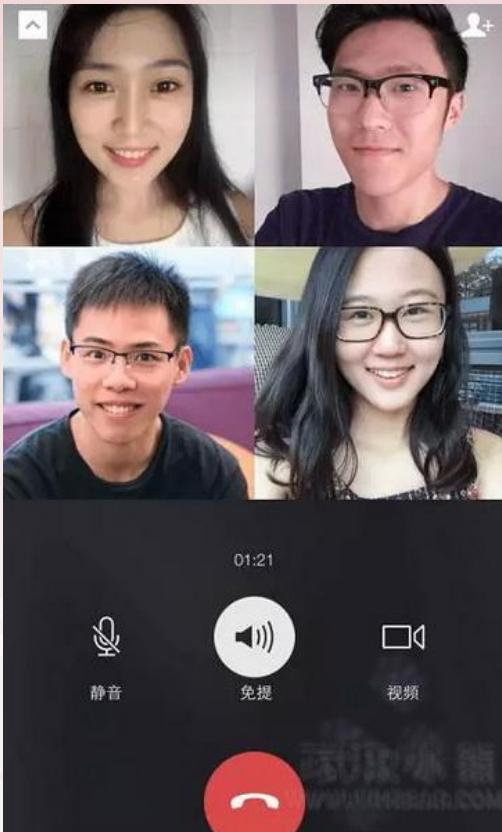


# 浅析低延迟直播协议设计： RTP/RTCP

# 哪些场景延迟必须要低？

## 微信群视频



## ME



# 为什么基于TCP的协议延迟不够低

## 无丢包

- 数据队列1-10
- 每秒传输1个单位数据
- 第10个单位的数据在第9s开始传输

## 有丢包

- 数据队列1-10
- 每秒传输1个单位数据
- 中间丢失了第4个数据，重传
- 第10个单位的数据在第10s开始传输

# 对丢包的处理是核心区别



搬砖的！嘿！叫你  
呢！把砖搬进去，一  
个都不能少！



把砖搬进去，自己看  
着办吧。



8砖/min  
掉砖捡起来+5s/砖

老板：为什么只有一  
块！？  
民工：前面的都掉在  
路上了



# Catalog

- RTP
- RTCP
- FEC
  - XOR
  - Reed Solomon
- DCCP
- A Custom Demo
- References

# RTP

## include

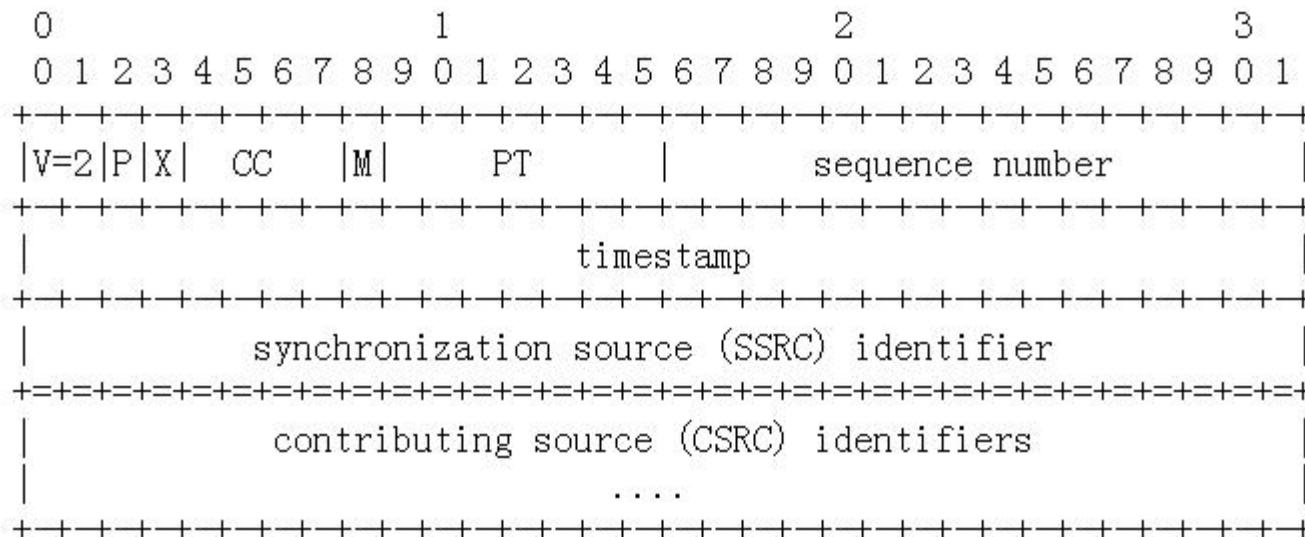
- end-to-end delivery services for data with real-time characteristics
- payload type identification
- sequence numbering
- Timestamping and delivery monitoring

## exclude

- ensure timely delivery or provide other quality-of-service guarantees
- delivery or prevent out-of-order delivery

# RTP Fixed Header Fields

The RTP header has the following format:



# Sample RTP network

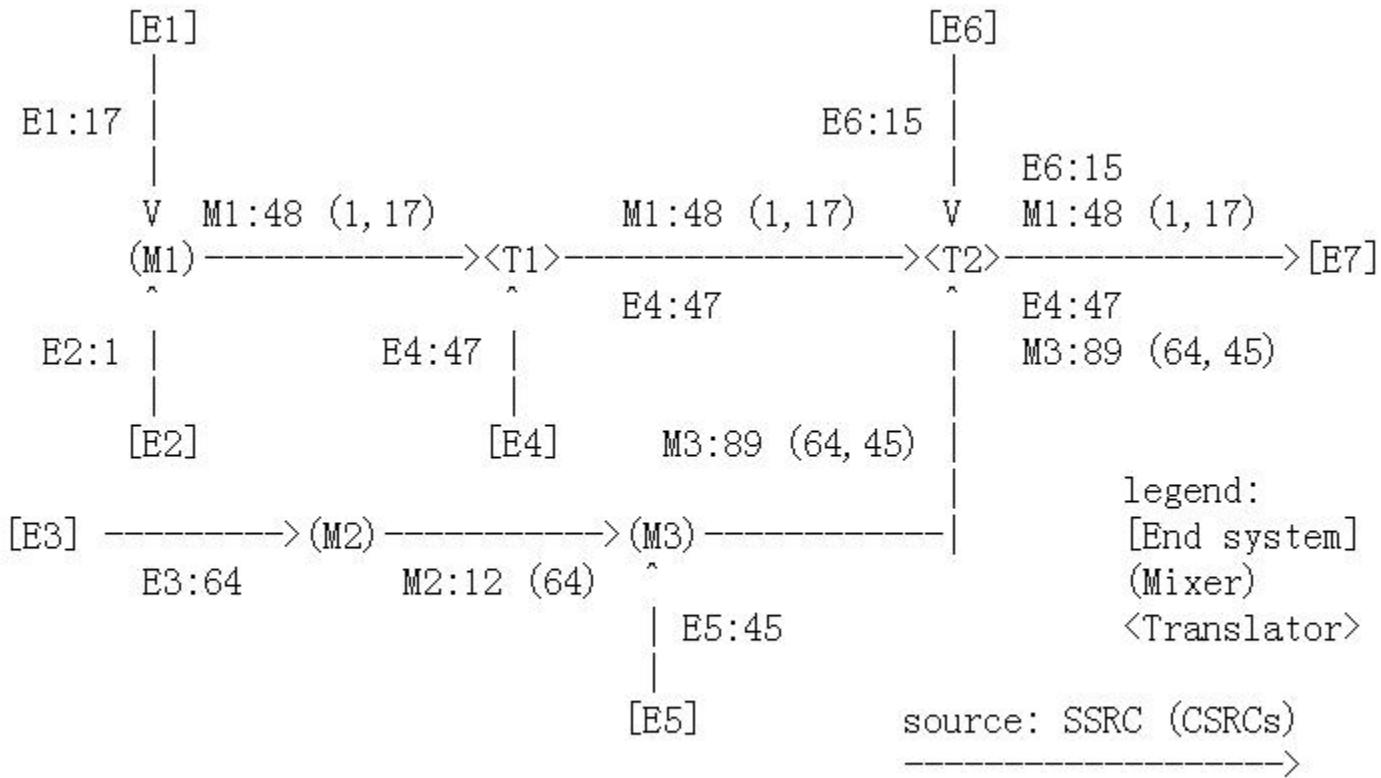


Figure 3: Sample RTP network with end systems, mixers and translators

# RTCP

- Feedback on the quality of the data distribution
- Persistent transport-level identifier
- Controlled rate to scale up
- Convey minimal session control information

# RTCP Sender Report

	0	1	2	3
header	V=2 P	RC	PT=SR=200	length
			SSRC of sender	
sender		NTP timestamp, most significant word		
info		NTP timestamp, least significant word		
		RTP timestamp		
		sender's packet count		
		sender's octet count		
report		SSRC_1 (SSRC of first source)		
block				
1	fraction lost	cumulative number of packets lost		
		extended highest sequence number received		
		interarrival jitter		
		last SR (LSR)		
		delay since last SR (DLSR)		



# RTCP Receiver Report

	0	1	2	3
header	0 1 2 3 4 5 6 7 8 9 0	1 2 3 4 5 6 7 8 9 0	1 2 3 4 5 6 7 8 9 0	1
report	V=2 P  RC   PT=RR=201   length			
block		SSRC of packet sender		
report		SSRC_1 (SSRC of first source)		
block				
1	fraction lost   cumulative number of packets lost			
		extended highest sequence number received		
		interarrival jitter		
		last SR (LSR)		
		delay since last SR (DLSR)		
report		SSRC_2 (SSRC of second source)		
block				
2	:	...	:	
		profile-specific extensions		



# Analyzing

## Interarrival Jitter

If  $S_i$  is the RTP timestamp from packet  $i$ , and  $R_i$  is the time of arrival in RTP timestamp units for packet  $i$ , then for two packets  $i$  and  $j$ ,  $D$  may be expressed as

$$D(i, j) = (R_j - R_i) - (S_j - S_i) = (R_j - S_j) - (R_i - S_i)$$

The interarrival jitter SHOULD be calculated continuously as each data packet  $i$  is received from source  $SSRC_n$ , using this difference  $D$  for that packet and the previous packet  $i-1$  in order of arrival (not necessarily in sequence), according to the formula

$$J(i) = J(i-1) + (|D(i-1, i)| - J(i-1))/16$$

Whenever a reception report is issued, the current value of  $J$  is sampled.

The interarrival jitter  $J$  is defined to be the mean deviation (smoothed absolute value) of the difference  $D$  in packet spacing at the receiver compared to the sender for a pair of packets.

## Round-Trip Time

```
[10 Nov 1995 11:33:25.125 UTC]          [10 Nov 1995 11:33:36.5 UTC]
n           SR(n)                         A=b710:8000 (46864.500 s)
----->
v           ^                                ^
ntp_sec =0xb44db705 v      dlsr=0x0005:4000 ( 5.250s)
ntp_frac=0x20000000 v      ^ lsr =0xb705:2000 (46853.125s)
(3024992005.125 s) v      ^ RR(n)
r           v                                ^
----->
|<-DLSR->|                                |
(5.250 s)                                     |
```

---

```
A     0xb710:8000 (46864.500 s)
DLSR -0x0005:4000 ( 5.250 s)
LSR  -0xb705:2000 (46853.125 s)
-----
```

---

```
delay 0x0006:2000 ( 6.125 s)
```

Figure 2: Example for round-trip time computation

# Options for Repair of Streaming Media

- Retransmission
- Forward Error Correction
  - Media-Independent FEC
  - Media-Specific FEC
- Interleaving



# Forward Error Correction

- FEC is one of the main methods used to protect against packet loss over packet-switched networks.



# Unequal Level Protection

Packet A	#####	#####
	:	:
Packet B	#####	#####
	:	:
ULP FEC Packet #1	00000000	00000000
	:	:
Packet C	#####	#####
	:	:
Packet D	#####	#####
	:	:
ULP FEC Packet #2	0000000000000000	0000000000000000
	:	:
	:<-L0->:	<--L1-->:

Figure 1: Unequal Level Protection



# FEC Packet

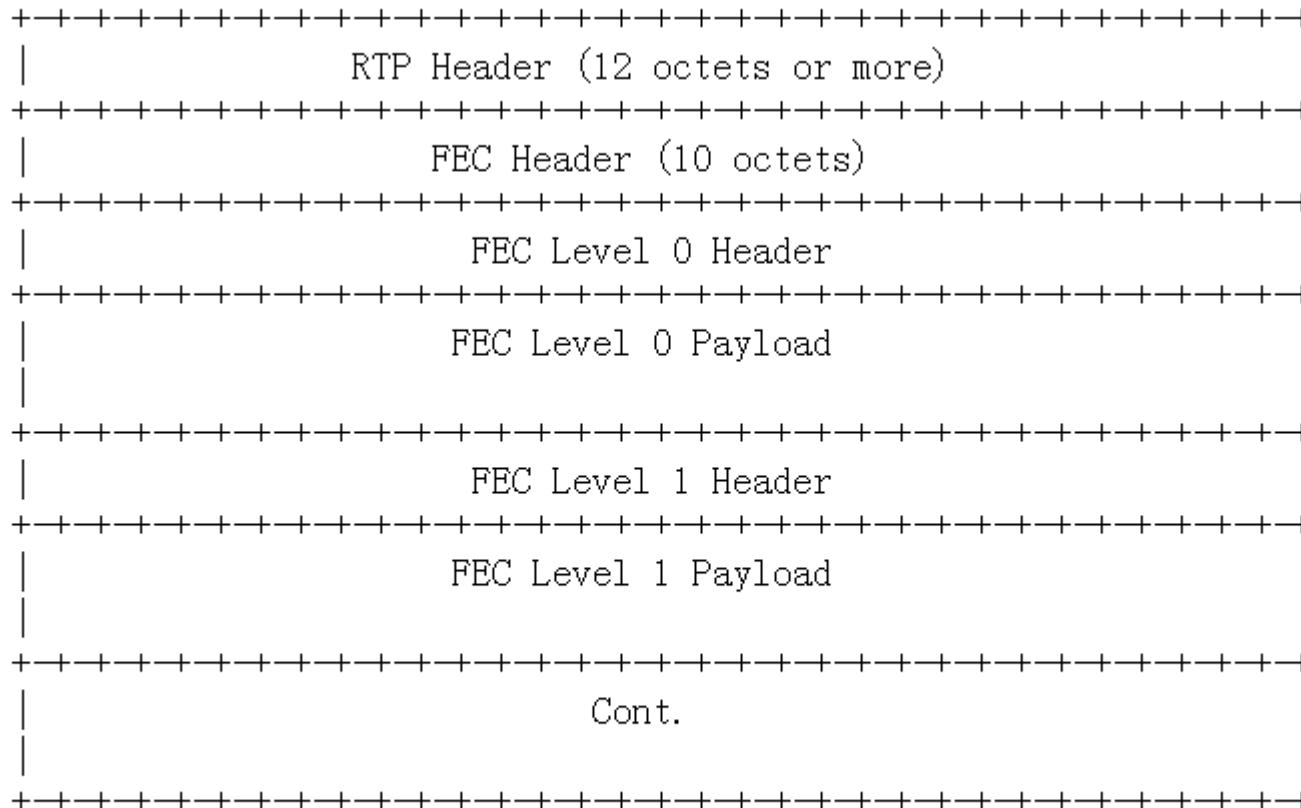


Figure 2: FEC Packet Structure

# XOR

- Data a,b,c,d
- Let e = XOR(a,b,c,d)
- Then
  - a = XOR(b,c,d,e)
  - b = XOR(a,c,d,e)
  - c = XOR(a,b,d,e)
  - d = XOR(a,b,c,e)

# Reed Solomon

$$\begin{array}{c} n \\ \overbrace{\quad\quad\quad\quad\quad} \\ \left( \begin{array}{ccccc} 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ B_{11} & B_{12} & B_{13} & B_{14} & B_{15} \\ B_{21} & B_{22} & B_{23} & B_{24} & B_{25} \\ B_{31} & B_{32} & B_{33} & B_{34} & B_{35} \end{array} \right) \\ \left\{ \begin{array}{l} n+m \\ B \end{array} \right\} \end{array} * \left( \begin{array}{c} D_1 \\ D_2 \\ D_3 \\ D_4 \\ D_5 \end{array} \right) = \left( \begin{array}{c} D_1 \\ D_2 \\ D_3 \\ D_4 \\ D_5 \\ C_1 \\ C_2 \\ C_3 \end{array} \right) \left\{ \begin{array}{l} D \\ C \end{array} \right\}$$

# Datagram Congestion Control Protocol

- The Datagram Congestion Control Protocol (DCCP) is a transport protocol that provides bidirectional unicast connections of congestion-controlled unreliable datagrams. DCCP is suitable for applications that transfer fairly large amounts of data and that can benefit from control over the tradeoff between timeliness and reliability.

# Brief demo:zoom



# Thanks



# References

- RFC3550
  - RTP: A Transport Protocol for Real-Time Applications
- RFC3551
  - RTP Profile for Audio and Video Conferences with Minimal Control
- RFC5109
  - RTP Payload Format for Generic Forward Error Correction
- RFC4340
  - Datagram Congestion Control Protocol (DCCP)
- RFC2354
  - Options for Repair of Streaming Media